

# SHA-3, Round 2: Hardware Implementations

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# Overview

- SHA-3 competition, round 2
- The candidates from a HW perspective
- An attempt at a fair comparison of existing HW implementations
  - High-speed implementations
  - Low-area implementations

# SHA-3 Competition

- Round 2: 50 down, 14 left
- Tweaks allowed till Sept. 15, 2009
- One year allocated for public comments
- NIST: “Hardware implementation will be addressed more thoroughly during the Round 2 analysis”
- HW implementations collected at ECRYPT2 “SHA-3 Zoo”



**I WANT YOUR HASH  
FOR SHA-3**

# “Still Alive”

- BLAKE
- Blue Midnight Wish
- CubeHash
- ECHO
- Fugue
- Grøstl
- Hamsi
- JH
- Keccak
- Luffa
- Shabal
- SHAvite-3
- SIMD
- Skein



(c) Valve

# The Candidates from a HW Perspective

- Concentrate on variants with 256-bit MD
- Neglect small details and fancy stuff, e.g. salting
- Main contributors to area
  - Storage
  - Compression function, message expansion function
- Functions on “critical path”
  - Dataflow within hash function

# BLAKE(-32)

- Local wide-pipe design
- Storage
  - Message block: 512 bits
  - Chaining value: 256 bits
  - Inner state of compression function: 512 bits
- Compression based on G function
  - Addition mod  $2^{32}$ , XOR, fixed rotation
  - 10 rounds à 8 G function invocations
  - Four G functions in parallel

# Blue Midnight Wish(-256)

- Wide-pipe design
- Storage
  - Message block: 512 bits
  - Chaining value: 512 bits
  - Inner state (for result of  $f_0$ ): 512 bits
- Compression based on  $f_0$ ,  $f_1$ ,  $f_2$  functions
  - Addition/subtraction mod  $2^{32}$ , XOR, fixed shift/rotation
  - $f_0$ : 16x 32-bit words in parallel
  - $f_1$ : 16x 32-bit words in sequence
  - $f_2$ : 16x 32-bit words in parallel
  - One invocation of  $f_0$ ,  $f_1$ ,  $f_2$  per message block

03/14

# CubeHash

- Uniform structure with high configurability
  - Message block: 8 – 1024 bits
- Storage
  - Chaining value: 1024 bit (32x 32-bit words)
- Compression based on simple rounds
  - Addition mod  $2^{32}$ , XOR, fixed rotation, 32-bit word swap
  - Number of rounds depends on variant (1 to  $\infty$ )

04/14

# ECHO(-256)

- Merkle-Damgård structure
- Message block: 1536 bits
- Storage
  - Chaining value: 512 bits
  - Inner state: 2048 bits (4 x 4 matrix of “AES states”)
- Compression based on AES transformations
  - 8 iterations of BIG.round:
    - BIG.SubWords: 16x 2 AES rounds in parallel
    - BIG.ShiftRows: Shifting of the matrix of “AES states”
    - BIG.MixColumns: 64x AES MixColumns in parallel

05/14

# Fugue(-256)

- Based on Grindahl
  - Message block: 32 bits
- Storage
  - Internal state/Chaining value: 960 bits (4 x 30 bytes)
- Compression based on AES-like operations
  - 2 sub-rounds per message block:
    - Rotation & XOR of 32-bit words
    - 16x AES S-box in parallel
    - 1x Super-Mix: 16 x 16 matrix multiplication
  - Rather expensive output transformation

# Grøstl(-256)

- Based on 2 parallel AES-like permutations
- Storage (each 8 x 8 bytes)
  - Message block: 512 bits
  - Chaining value: 512 bits
  - Intermediate state: 512 bits
- Compression with P & Q permutations
  - 10 rounds à 4 transformations:
    - AddRoundConstant
    - SubBytes: 64 AES S-boxes in parallel
    - ShiftBytes: Rotate 8-byte rows
    - MixBytes: 8 x 8 matrix multiplication

07/14

# Hamsi(-256)

- “Concatenate-Permute-Truncate” design
- Message block: 32 bits
  - Expanded to 256 bits by linear code
- Storage
  - Chaining value: 256 bits
  - Internal state of permutation: 512 bits
- Non-linear permutation P / Pf
  - P applied 3 times per message block:
    - Addition of counter & constant
    - 128 (4 x 4) S-boxes in parallel
    - 4x linear transformation L in parallel ( $2^{128} \rightarrow 2^{128}$ )
  - Output transformation: Pf applied 6 times on last block

## JH(-256)

- Based on balanced Feistel network
- Storage
  - Message block: 512 bits
  - Internal state/Chaining value: 1024 bits
  - Round constant: 256 bits
- Compression based on round function  $R_8$ 
  - $R_8$  applied 35 times on internal state & round constant
    - 256 & 64 (4 x 4) S-boxes in parallel
    - 128 & 32 linear transforms L over GF( $2^4$ ):  $2^4 \rightarrow 2^4$
    - Permutation of 256x & 64x 4-bit words

09/14

# Keccak(-256)

- Sponge function construction
- Message block: 1024 bits
- Storage
  - Internal state: 1600 bits (25x 64-bit words)
- **Round R, iterated 18 times per message block**
  - XOR, AND, NOT
  - 25x bit permutation within words in parallel
  - Permutation of words
  - Addition of round constants

# Luffa(-256)

- Based on round with message injection & permutation
- Message block: 256 bits
- Storage
  - Chaining value: 768 bits (3x 256-bit blocks)
- Message injection
  - Multiplication by constants in  $GF(2^8)^{32}$  of 256-bit blocks (chaining value, message block)
  - Addition (XOR) of results
- Permutation: 3  $Q_j$  (256-bit permutation) in parallel
  - 8 iterations of a step function
    - SubCrumb: 192 (4 x 4) S-boxes in parallel
    - MixWords: 12 linear 64-bit permutations (fixed rotations, XORs) in parallel
- Output transformation uses round with all-zero message

# Shabal(-256)

- NLFSR-based compression function
- Storage
  - Message block: 512 bits
  - Chaining value: 1408 bits
  - Message block counter: 64 bits
- Compression function  $f_{Sh}$ 
  - Addition/subtraction mod  $2^{32}$ , XOR
  - 896-bit permutation P
    - 16 fixed rotations of 32-bit words in parallel
    - 3x 16 shifts of the NLFSR structure in sequence
      - Constant multiplication mod  $2^{32}$ , fixed rotations of 32-bit words, XOR, AND, NOT
    - 3x 12 parallel additions mod  $2^{32}$  in sequence
  - Output transformation: Invokes  $f_{Sh}$  3 times

# SHAvite-3(<sub>256</sub>)

- Based on HAIFA and AES
- Storage
  - Message block: 512 bits
  - Chaining value: 256 bits
  - Bit counter: 64 bits
- Compression function based on balanced Feistel network ( $E^{256}$  block cipher)
  - Function f: 3 AES rounds
  - 12 rounds
- Message expansion produces AES round keys
  - Non-linear step: 4 AES rounds, XOR with bit counter
  - Linear step: XOR previous round keys

# SIMD(-256)

- FFT-based message expansion
- Storage
  - Message block: 512 bits
  - Chaining value: 512 bits
  - Intermediate state: 512 bits
- Message expansion  $2^{512} \rightarrow 2^{4096}$  (3 layers)
  - Number-theoretic transform (FFT)
  - Inner code
  - Permutation
- Compression function
  - Uses block cipher with 4 parallel Feistel ladders (128 bits each)
    - Addition mod  $2^{32}$ , step-dependent rotation, AND, OR, NOT
  - 32 iterations in total (128 bits of exp. message per step)
  - 4 extra steps at end

# Skein(-x-256)

- Based on tweakable block cipher Threefish
- Storage
  - Message block: 256 / 512 / 1024 bits
  - Chaining value: 256 / 512 / 1024 bits
  - Intermediate state: 256 / 512 / 1024 bits
  - Expanded key & tweak: 512 / 768 / 1280 bits
- Compression / Output transf. with Threefish
  - 72 / 72 / 80 rounds
    - 2 / 4 / 8 additions mod  $2^{64}$  in parallel
    - Round-dependent rotation, XOR, permutation of 64-bit words
    - Key expansion: 3 additions mod  $2^{64}$  in parallel

# Summary of Features

Candidate	Storage bits	„Throughput-relevant“ features
BLAKE-32	1,280	10 rounds: 2x G function in sequence; Add mod $2^{32}$
BMW-256	1,536	Only one $f_0$ , $f_1$ , $f_2$ per message block; Add/sub mod $2^{32}$
CubeHash	1,024	Number of rounds; Size of message block; Add mod $2^{32}$
ECHO-256	2,560	8 iterations of BIG.round
Fugue-256	960	32-bit message block; 2 sub-rounds; AES S-box & Super-Mix
Grøstl-256	1,536	10 rounds; AES S-box & MixBytes
Hamsi-256	768	3x P permutation; (4 x 4) S-box; linear transform
JH-256	2,048	35 iterations of R8; (4 x 4) S-box; linear transform
Keccak-256	1,600	18x round R; Boolean functions
Luffa-256	768	8 iterations of step function; (4 x 4) S-box; linear permutation
Shabal-256	1,984	Add mod $2^{32}$ ; 3x 16 shifts & 12 additions in sequence
SHAvite-3 <sub>256</sub>	832	12 rounds: 3x AES round; Message expansion: 4x AES round
SIMD-256	1,536	Number-theoretic transform (FFT) in message expansion
Skein-256-256	1,280	72 rounds; add mod $2^{64}$
Skein-512-256	2,304	72 rounds; add mod $2^{64}$
Skein-1024-256	4,352	80 rounds; add mod $2^{64}$

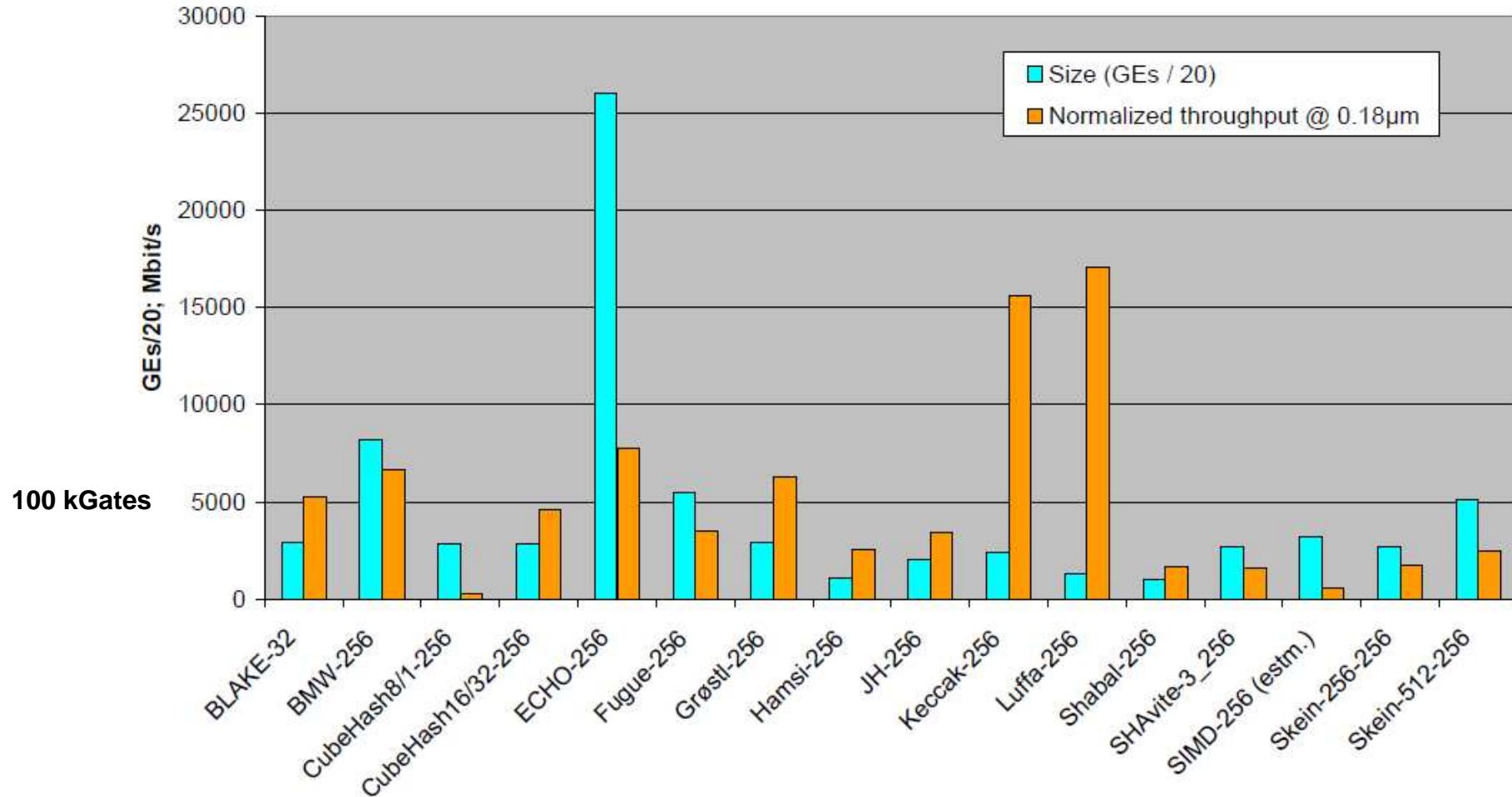
# To Tweak or Not to Tweak ... (updated)

Candidate	Round 2 tweaks (tweaks due Sept. 15, 2009; status as of Sept. 17, 7:30 CEST)
BLAKE	None
BMW	Changed inputs to $f_0$ and $f_1$ ; Additional invocation of compression function at end
CubeHash	2x rounds ( $r = 16$ ); 32x size of message block ( $b = 32$ )
ECHO	None
Fugue	Not yet known
Grøstl	None
Hamsi	None (but additional variants specified)
JH	None
Keccak	Not yet known
Luffa	Modification of S-box; Changed order of SubCrumb inputs; Always one round in output transformation
Shabal	Not yet known
SHAvite-3	Inversion of some counter values
SIMD	Not yet known
Skein	Changed rotation constants

# High-Speed HW Implementations

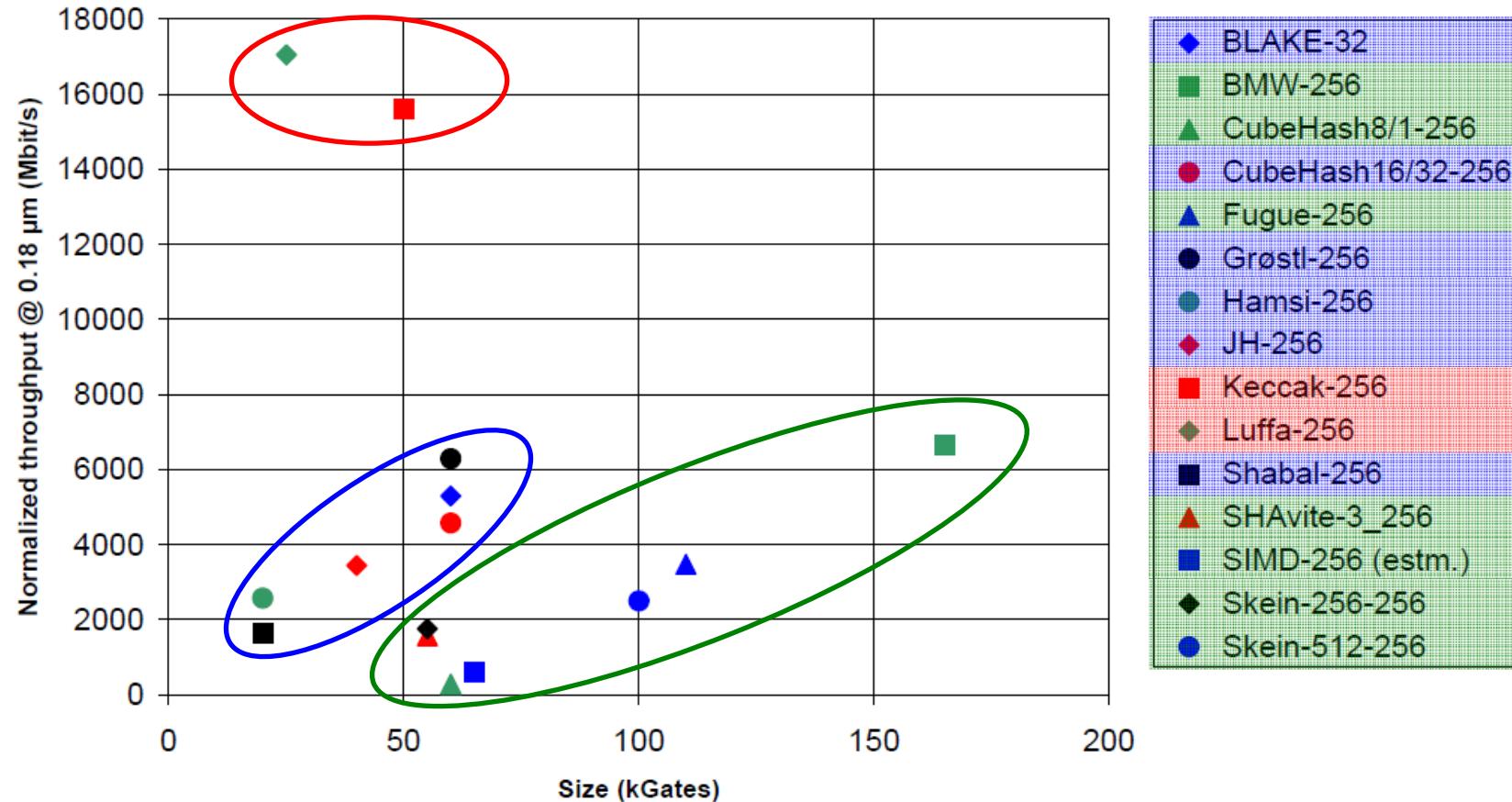
- Standard-cell implementations
  - Primary goal: High throughput
  - Secondary goal: Low area
- Direct comparison is problematic! Differences in
  - Implementation scope (full, ext. memory, core)
  - HW module interface
  - Process technology
  - Standard-cell library
  - Implementer's design experience & effort
  - Synthesis effort
  - ...
- Using **very rough** scaling to uniform minimal gate length (0.18 µm)

# High-Speed HW Implementations



- Throughput scaled with square of the minimal gate length to 0.18 μm
- Results for CubeHash, JH, SHAvite-3, SIMD preliminary and pending publication
- BLAKE, BMW, and Shabal: Core functionality only

# High-Speed HW Implementations

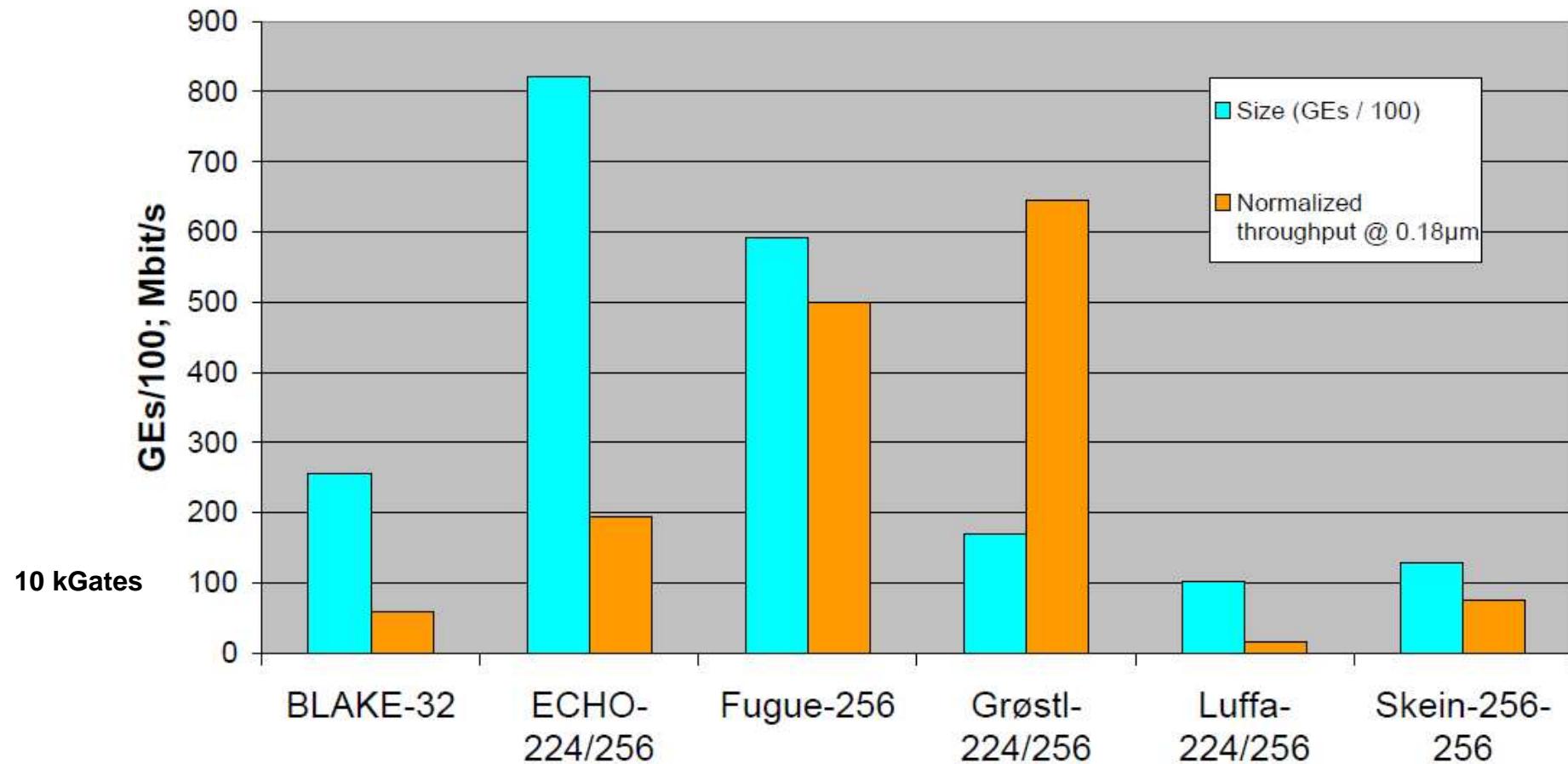


- Throughput scaled with square of the minimal gate length to 0.18 μ
- ECHO-256 omitted: ~ 7750 Mbit/s with 520 kGates
- Results for CubeHash, JH, SHAvite-3, SIMD preliminary and pending publication
- BLAKE, BMW, and Shabal: Core functionality only

# Low-Area HW Implementations

- Primary target: Minimal area and/or power/energy consumption
- Raw throughput of lesser importance
- Comparison only of fully autonomous implementations
- Only 6 candidate implementations available so far
- To be taken with a grain of salt!
  - Some implementations not totally optimized towards low-area

# Low-Area HW Implementations



- Throughput scaled with square of the minimal gate length to 0.18 μm
- Comparison includes only fully autonomous implementations

# All the Figures ...

- ... on the pages of the ECRYPT2 “SHA-3 Zoo”:  
[http://ehash.iaik.tugraz.at/wiki/The\\_SHA-3\\_Zoo](http://ehash.iaik.tugraz.at/wiki/The_SHA-3_Zoo)
- All reported ASIC & FPGA implementations
- Frequent updates for the candidates

# Thanks for your attention!

Questions?



Shoot!